

CS 4350: Fundamentals of Software Engineering
CS 5500: Foundations of Software Engineering

Lesson 8.1 Static Program Analysis

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Alternatives to Testing

- So “program testing can be used to show the presence of bugs, but never to show their absence” (Dijkstra’s Law), what can we do?
- Testing is limited to finite concrete cases;
 - Can we check unbounded symbolic cases?
- Yes! *

*Some restrictions apply:

- Can show absence, but cannot show presence;
- Sometimes cannot show either;
- How much time do you have?
- ...

Outline of this lesson

1. The impractical goal of **program verification**.
2. What lies between testing and verification?
 - a. Partial verification
 - b. Optional type systems
 - (Should be familiar: TypeScript anyone?)
 - c. Bug finders
 - (Also familiar: es-lint)

Learning Objectives for this Lesson

- By the end of this lesson, you should be able to:
 - List challenges preventing full verification;
 - Define false/true negatives/positives;
 - Variously describe the effects of false positives and negatives;
 - Explain when to integrate static analysis into builds.

Verification Checks Code Against Specification



A Specification: What the system is supposed to do.

B Code: What the system actually does.

- Difficulties:
 - Need a **full formal** specification;
 - Even proving termination is undecidable, let alone proving adherence to a specification.

A Full Formal Specification

- ... precisely defines exactly the behavior that the system should have:
 - What the outputs are in terms of the inputs;
 - What behaviors the system should have.
- Wait a minute! That's called a program.
- Yes, a full formal specification is essentially a program, perhaps expressed at a **higher level.**
 - ... with all the complexity that entails,
 - ... including bugs!

Inefficiently
(w/o algorithms)

We can't avoid
Human fallibility.

Specification Must Be Modular

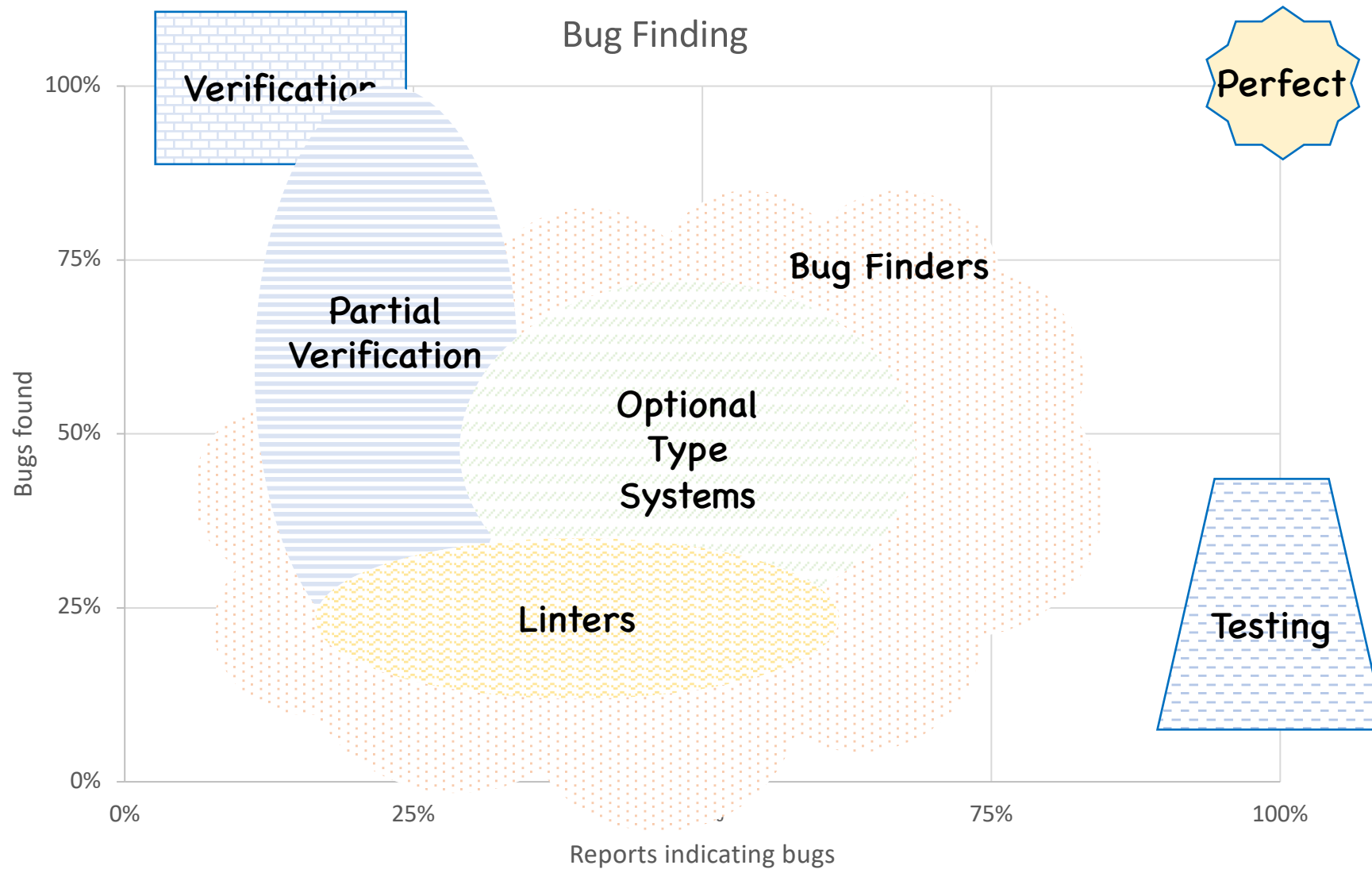
- Without modularity
 - Specification is incomprehensible
 - It is likely *inadequate* (i.e., doesn't specify what we want).
 - Specification is unprovable
 - Proof checking is usually exponential or worse
 - Must break down into usable pieces.
- Specification has to be maintained as code is:
 - Every function/class/module needs specification.
 - Even every loop needs its own specification.

Proof Must Be Modular Too

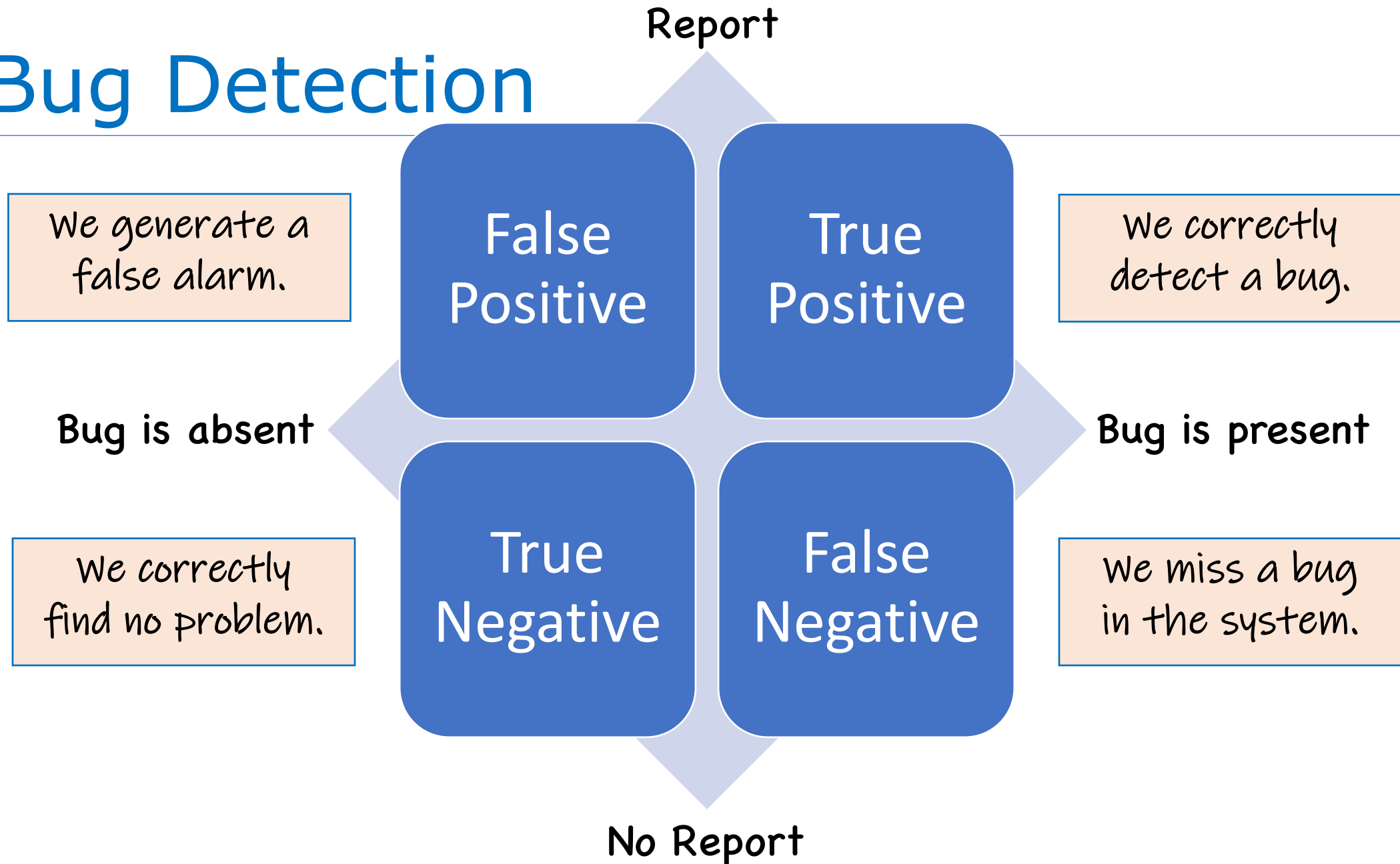
- A verification proof is usually very complex, needing lemmas written by hand.
 - Typically written and stored along with the specification.
- Engineering proofs is a highly specialized skill:
 - Hint: harder than coding → More \$\$\$
 - Proof must be updated
 - **every time** program or specification changes!
- Usually too expensive unless safety critical **and** mandated by regulation.

Verification Doesn't Prove Presence of Bugs

- Verification fails if
 - Missing lemma for unit behavior, or
 - Cannot verify loop invariant, or
 - Functional specification missing a piece, or
 - Run out of time trying to construct proof, or
 - Specification is wrong.
- Constructing the proof can easily take as long as constructing the software, if not much more.
 - Just because there is no proof does not mean the software has a fault.



Bug Detection



Static Program Analysis

- Bug finders and partial verification use static program analysis
 - Reads in the program (just like a compiler);
 - Analyze to determine properties:
 - E.g., are all open resources eventually closed?
 - "static" means "without running the program".
- All non-trivial properties are undecidable
 - Approximations are always necessary: make a choice
 - E.g., miss some closes of open resources, or
 - Miss some open resources not being closed.

Compromises with Static Analysis

- Getting precise results may take **time**:
 - Many algorithms are exponential in precision measures.
- Getting precise results may require **whole** program:
 - If parts of the program loaded at runtime:
 - Analysis results may be very imprecise, or (worse)
 - Incorrect, if they assume the whole program is available.
- Getting precise results may require intervention:
 - Code may need to be **annotated** with information:
 - E.g., this method may return an open resource.

Effects of Analysis Imprecision

FALSE NEGATIVES

- The static analysis misses something "bad" in program:
 - Bug not found.
- Can give a false sense of security.
- Can be reduced, but at the cost of false positives!

FALSE POSITIVES

- The static analysis reports a problem that doesn't exist:
 - There is no bug.
- Real bugs can be swamped by a flood of spurious reports.
- Programmer time is wasted chasing down false leads.

Google defined “Effective False Positive”

- A report from static analysis is effectively false,
 - If it is ignored by developers;
 - Whether or not it represents a true bug.
- Even if the report is technically correct
 - It may refer to something considered unimportant:
 - E.g., who cares if all the files aren't closed, if the program is about to exit anyway.
 - E.g., yes, there is a race condition between two logging statements, but that's not important.
- Even if the report is technically wrong
 - Developers may see potential problem, and fix.

Criteria For Automated Program Analysis

- Efficient and Easy
 - Should not require whole program or annotations.
- Rarely spurious
 - No more than 10% effectively false positive.
- Actionable
 - Should point out things easy to fix.
- Effective
 - Problems should be perceived as important.

*Automatically applied
during Code Review.*

Source: Software Engineering at Google, [Chapter 20](#)

Review: Learning Objectives for this Lesson

- You should now be able to:
 - List challenges preventing full verification;
 - Define false/true negatives/positives;
 - Variously describe the effects of false positives and negatives;
 - Explain when to integrate static analysis into builds.

Next steps...

- In our next lesson, we'll talk about Code Smells and Refactoring.